

Solutions to Homework 5 - Markov Chains

1) Stationary distribution. Consider a Markov chain (MC) $X_N = X_0; X_1; \dots; X_n; \dots$ with state space $S = \{1; 2\}$ and transition probability matrix

$$P = \begin{pmatrix} 1/4 & 3/4 \\ 1/5 & 4/5 \end{pmatrix}$$

To obtain the stationary distribution $\pi = [\pi_1; \pi_2]$

$$\lim_{n \rightarrow \infty} P_{22}^n = \pi_2 = \frac{15}{19}$$

Finally, by virtue of the ergodic theorem the time average $\frac{1}{n} \sum_{k=0}^{n-1} \mathbb{1}_{\{X_k = 1\}}$ also converges. The long-run fraction of time the MC visits state 1 is thus given by

$$\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{k=0}^{n-1} \mathbb{1}_{\{X_k = 1\}} = \pi_1 = \frac{4}{19}$$

2) A cloudy town. A certain town never has two sunny days in a row. Each day is classified as being either sunny, cloudy (but dry), or rainy. If it is sunny one day, then it is equally likely to be either cloudy or rainy the next day. If it is rainy or cloudy one day, then there is one chance in two that it will be the same the next day, and if it changes then it is equally likely to be either of the other two possibilities.

Let $X_N = X_0; X_1; \dots; X_n; \dots$ be the random process describing the weather evolution of the given town, with n denoting the day number. Given the nature of the evolution of the process (meaning the weather today only

Customers who do not find the product in stock depart without making a purchase. The store orders new units of the product from its supplier at the end of the day (after that day's demand has materialized). However, the supplier is not completely reliable, and each day, with probability p dependent of everything else, the order is permanently lost in which case the order does not arrive to the store. If the order is not lost, it arrives to the

Now, what is left is to carefully put all pieces (2)-(6) together back in (1). For all i, j , the transition probabilities are thus given by

$$\begin{aligned}
 P_{i0} &= \sum_{k=i}^{\infty} p(k); \\
 P_{ij} &= \begin{cases} p(i-j); & i \leq j \\ 0; & i > j \end{cases}; \quad \text{for } 0 < j < q; \\
 P_{iq} &= \begin{cases} p(i-q) + \sum_{k=i}^{q-1} p(k); & i < q \\ \sum_{k=i}^{q-1} p(k); & i \geq q \end{cases}; \\
 P_{ij} &= \begin{cases} p(i-j) + \sum_{k=i}^{q-1} p(k); & i \leq j < q \\ \sum_{k=i}^{q-1} p(k); & i > j < q \\ 0; & i \leq j < q \text{ and } i > j + q \\ 0; & i > j + q \text{ and } i > j < q \end{cases}; \quad \text{for } j > q;
 \end{aligned}$$

4) Non-invertible function of a Markov chain. Suppose that $X_N = X_0; X_1; \dots; X_n; \dots$ is a MC with state space $S = \{1; 2; 3\}$, transition probability matrix

$$P = \begin{pmatrix} 0 & 0 & 1 \\ 1/4 & 2/3 & 1/3 \\ 3/4 & 1/4 & 0 \end{pmatrix}$$

and initial distribution $P[X_0 = 1] = 1/5$, $P[X_0 = 2] = 2/5$ and $P[X_0 = 3] = 2/5$. Suppose that the random process $Y_N = Y_0; Y_1; \dots; Y_n; \dots$ satisfies $Y_n = g(X_n)$, $n \geq 0$, where $g(1) = 1$ and $g(2) = g(3) = 2$.

To calculate $P\{Y_2 = 1, Y_1 = 2; Y_0 = 1\}$, one can resort to the definition of conditional probability and write

$$\begin{aligned}
 P\{Y_2 = 1, Y_1 = 2; Y_0 = 1\} &= \frac{P\{Y_2 = 1; Y_1 = 2; Y_0 = 1\}}{P\{Y_1 = 2; Y_0 = 1\}} \\
 &= \frac{P\{X_2 = 1; X_1 = 2; X_0 = 1\}}{P\{X_1 = 2; X_0 = 1\}} \\
 &= \frac{P\{X_2 = 1; X_1 = 2; X_0 = 1\} + P\{X_2 = 1; X_1 = 3; X_0 = 1\}}{P\{X_1 = 2; X_0 = 1\} + P\{X_1 = 3; X_0 = 1\}} \\
 &= \frac{P\{X_0 = 1\} P_{12} P_{21} + P\{X_0 = 1\} P_{13} P_{31}}{P\{X_0 = 1\} P_{12} + P\{X_0 = 1\} P_{13}} = \frac{5}{12}.
 \end{aligned}$$

Because the function g is not invertible, it turns out that Y_N is not Markov chain. Actually, one can verify that

$$P\{Y_2 = 1, Y_1 = 2; Y_0 = 1\} \neq P\{Y_2 = 1, Y_1 = 2\} = \frac{17}{36}.$$

5) A non-irreducible Markov chain. Consider a MC $X_N = X_0; X_1; \dots; X_n; \dots$ with state space $S = \{1; 2; 3\}$ and transition probability matrix

$$P = \begin{pmatrix} 0 & 1/3 & 1/6 & 1/2 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

A) The

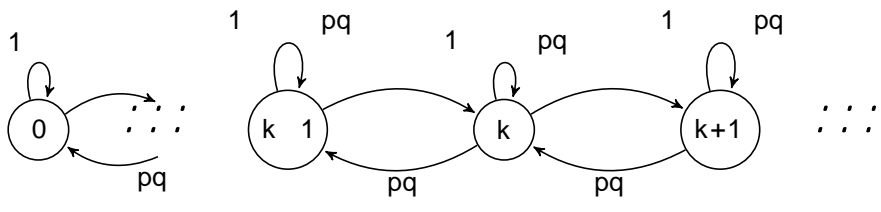


TABLE I
SUMMARY OF THE RECURRENCE ANALYSIS RECURRENCE DEPENDS ON THE PARAMETERS AND pq .

Parameters	Effect on recurrence
$= 0, pq = 0$	Every state is a class. Every class is positive recurrent.
$= 0, pq > 0$	Every state is a class. Class 0 is positive recurrent. All other classes are transient.
$> 0, pq = 0$	Every state is a class. Every class is transient.
$> pq > 0$	The MC is irreducible and transient.
$pq > > 0$	ys

In order to obtain the value of the constant, we use the fact that the limit probabilities π_k of every state $k = 0; 1; 2; \dots$ in the MC must add up to 1. Formally,

$$1 = \sum_{k=0}^{\infty} \pi_k = c \sum_{k=0}^{\infty} \left(\frac{p}{q}\right)^k = c \sum_{k=0}^{\infty} \left(\frac{p}{q}\right)^k = c \frac{1}{1 - \frac{p}{q}}; \quad (20)$$

where we used the fact that $\frac{p}{q} < 1$ for the convergence of the infinite geometric series. From (20) we obtain $c = 1 - \frac{p}{q}$ resulting in the limit probabilities

$$\pi_k = c \left(\frac{p}{q}\right)^k = \left(1 - \frac{p}{q}\right) \left(\frac{p}{q}\right)^k; \quad (21)$$

Observe that in (21), as $\frac{p}{q}$ tends to 1 from below, probabilities π_k tend to 0 for all k , which is consistent with the fact that when $\frac{p}{q} = 1$ the MC is null recurrent. Similarly, when $\frac{p}{q} > 1$ and the MC is transient, probabilities $\pi_k = 0$

which is true for $|x| < 1$ and that we used in (20). After differentiating both sides of (28) and multiplying both of them by x , we obtain

$$\sum_{k=0}^{\infty} kx^k = \frac{x}{(1-x)^2} \quad (29)$$

By using (29) for $x = \frac{\lambda}{\mu}$, we can evaluate the expected value expression in (27), to obtain

$$\lim_{n \rightarrow \infty} E[Q_n] = \frac{\lambda}{\mu} \frac{\frac{\lambda}{\mu}}{1 - \frac{\lambda}{\mu}} = \frac{\lambda}{\mu - \lambda} = \frac{\lambda}{\mu(1 - \frac{\lambda}{\mu})} \quad (30)$$

Notice that the expression obtained in (30) for the expected queue length is quite intuitive. The expected length grows with increasing arrival rate and decreases when the difference between the successful transmission rate and the arrival rate increases.

E) Probability of successful transmission and optimal transmission probability p . So far, we have assumed that the probability q that a transmission by an arbitrary terminal does not experience a collision with any other terminal is given. However, under the dominant system assumption (B), for a transmitted packet not to collide it must be that none of the remaining $J-1$ transmitted any packet. Since every terminal acts independently, we may write

$$q = (1-p)^{J-1} \quad (31)$$

This allows us to compute the probability that maximizes the probability of successful transmission. In order to do this, we differentiate pq with respect to p and look for the roots of the corresponding equation, i.e.

$$\frac{d}{dp} [pq] = \frac{d}{dp} p(1-p)^{J-1} = (1-p)^{J-1} - (J-1)p(1-p)^{J-2} = 0 \quad (32)$$

Since we obviously assume $p < 1$ because otherwise the probability of successful transmission is trivially null, we may divide (32) by $(1-p)^{J-2}$ to obtain a linear expression that yields the optimal probability

$$p = \frac{1}{J} \quad (33)$$

Since the probability of the queue being empty is an increasing function of p [cf. (22)] and the expected queue length is a decreasing function of p [cf. (30)], the optimal probability p in (33) entails simultaneously shorter queues and higher probabilities of these queues being empty. Moreover, for this probability of no collision becomes [cf. (31)]

$$q = 1 - \frac{1}{J} \quad (34)$$

If we consider a system with many terminals, we may estimate the probability of no collision as the limit of when J tends to infinity, i.e.,

$$\lim_{J \rightarrow \infty} q = \lim_{J \rightarrow \infty} \left(1 - \frac{1}{J}\right)^{J-1} = \lim_{J \rightarrow \infty} \left(1 - \frac{1}{J}\right)^J \left(1 - \frac{1}{J}\right)^{-1} = \frac{1}{e} \approx 0.368 \quad (35)$$

The above result implies that RA communications utilizes approximately 37% of the available access point resources without any coordination overhead among terminals.

F) Average time occupancies. It is possible to argue that the limit probabilities in (21) as well as the performance indicators in (22)-(24), and (27) are of little practical value. What these probabilities express is an average across all possible paths of the communication system. Say we run the system once and obtain a certain path

Fig. 4. Evolution of the first four queues over the first 1000 time slots. Queue 3 achieves a maximum queue length of seven packets. Every other queue length remained under this value.

these probabilities cannot be computed in closed form. However, we may use our simulation results to estimate the probability distribution function in (38). Moreover, since the MC is ergodic, we may compute the probabilities in (38) as the time averages over a single simulation run, that is

$$k \quad \frac{1}{N} \sum_{n=1}^N I f R_{j_n} = kg; \quad \text{for all } k; \quad (39)$$

and for large N . The Matlab script to generate the requested plot follows:

```
clear all; close all; clc;

J=16;
p=1/J;
N=10^5;
lambd=0.9*p*(1-p)^(J-1);

x = aloha_uplink_simulation(J, p, lambd, N);

Q1=max(x(1,:));
frequencies=zeros(1, Q1+1);
for i=0:Q1
    frequencies(1, i+1)=sum(x(1,:)==i);
end

rho=lambd/(p*(1-p)^(J-1));

A=[frequencies/N; (1-rho)*(rho.^(0:Q1))];
A=A';
figure
```

